


CMU SCS

**Carnegie Mellon Univ.
Dept. of Computer Science
15-415/615 - DB Applications**

C. Faloutsos – A. Pavlo
Lecture#1: Introduction




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Outline

- Introduction to DBMSs
- The Entity Relationship model
- The Relational Model
- SQL: the commercial query language
- DB design: FD, 3NF, BCNF
- indexing, q-opt
- concurrency control & recovery
- advanced topics (data mining, multimedia)

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We'll learn:

- What are RDBMS
 - when to use them
 - how to model data with them
 - how to store and retrieve information
 - how to search quickly for information
- Internals of an RDBMS: indexing, transactions

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We'll learn (cnt'd)

- Advanced topics
 - multimedia indexing (how to find similar, eg., images)
 - data mining (how to find patterns in data)

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Administrivia

- Weights: as announced

```

graph TD
    CG[Course grade] -- 30% --> ASGN[ASGN]
    CG -- 30% --> MT[MT]
    CG -- 40% --> FE[Final exam]
    ASGN -- 5% --> ASGN1[ASGN1]
    ASGN -- 5% --> ASGN8[ASGN8]
    ASGN1 -.- ASGN8
  
```

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Administrivia - II

- FYI: ASGN3 and ASGN7 are heavy
- Late policy: 4 'slip days'
- Exams: no aids allowed, except
 - 1 page with your notes (both sides) for MT
 - 2 such pages for Final

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Detailed outline

- Introduction
 - ➔ – Motivating example
 - How do DBMSs work? DDL, DML, views.
 - Fundamental concepts
 - DBMS users
 - Overall system architecture
 - Conclusions

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What is the goal of rel. DBMSs

(eg., you have 50 friends + phone#;
Or a dentist has 100 customers, addresses,
visit-info, treatment-info)
How can RDBMSs help?

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What is the goal of rel. DBMSs

Electronic record-keeping:
Fast and convenient access to information.

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Definitions

- ‘DBMS’ = ‘Data Base Management System’:
 the (commercial) system, like:
 DB2, Oracle, MS SQL-server, ...
- ‘Database system’: DBMS + data + application programs

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Motivating example

Eg.: students, taking classes, obtaining grades;

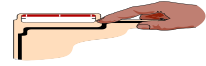
- find my gpa
- <and other ad-hoc queries>

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Obvious solution: paper-based

- advantages?
- disadvantages?



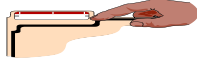
eg., student folders,
alpha sorted

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Obvious solution: paper-based

- advantages?
 - cheap; easy to use
- disadvantages?



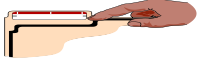
eg., student folders,
alpha sorted

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Obvious solution: paper-based

- advantages?
 - cheap; easy to use
- disadvantages?
 - no 'ad hoc' queries
 - no sharing
 - large physical foot-print




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Next obvious solution

- computer-based (flat) files +
- C (Java, ...) programs to access them



e.g., one (or more) UNIX/DOS files,
with student records and their courses


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Next obvious solution

your layout for the student records?

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Next obvious solution

your layout for the student records?
(eg., comma-separated values ‘csv’)

```
Smith,John,123,db,A,os,B
Tompson,Peter,234
Atkinson,Mary,345,os,B,graphics,A
```

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Next obvious solution

your layout for the student records?
(many other layouts are fine, eg.:

| | | |
|-------------------|--|----------------|
| Smith,John,123 | | 123,db,A |
| Tompson,Peter,234 | | 123,os,B |
| Atkinson,Mary,345 | | 345,os,B |
| | | 345,graphics,A |


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Problems?

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Problems?

- inconvenient access to data (need ‘C++’ expertise, plus **knowledge** of file-layout)
 - data isolation
- data redundancy (and inconsistencies)
- integrity problems
- atomicity problems

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Problems? (cont'd)

- ...
- concurrent-access anomalies
- security problems

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Problems? (cont'd)

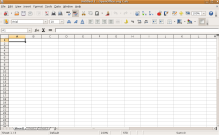

[why?

because of two main reasons:

- **file-layout** description is buried within the C programs and
- **Transactions**: there is no support for them (concurrency and recovery)

]

DBMSs handle exactly these two problems

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DBMS solution

- commercial/freeware DBMS &
- application programs

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Main vendors/products

| <u>Commercial</u> | <u>Open source</u> |
|-------------------|----------------------|
| • Oracle | Postgres (UCB) |
| • IBM/DB2 | mySQL/mariaDB |
| • MS SQL-server | sqlite (sqlite.org) |
| • Sybase | miniBase (Wisc) |
| • (MS Access, | Predator (Cornell) |
| • ...) | (www.acm.org/sigmod) |

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<Demo with sqlite3>

- Insert 'student' and 'takes' records
- Find the 'os' class roster
- Find the GPA of 'Smith'

www.cs.cmu.edu/~christos/courses/dbms.S15/files/sqldemo.zip

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How do DBs work?

Pictorially:

select *
from student

DBMS

data

and meta-data =
catalog =
data dictionary

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How do DBs work?

```
% sqlite3 miniu.sql
sqlite>create table student (
  ssn fixed;
  name char(20) );
```

| student | |
|---------|------|
| ssn | name |
| | |

Smith,John, 123,db,A,os,B
 Tompson,Peter,234
 Atkinson,Mary,345,os,B,graphics,A

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How do DBs work?

```
% sqlite3 miniu.sql
sqlite>create table student (
  ssn fixed;
  name char(20) );
```

| student | |
|---------|------|
| ssn | name |
| | |

| | |
|-----------|---------------------|
| Smith, | 123,db,A,os,B |
| Tompson, | 234 |
| Atkinson, | 345,os,B,graphics,A |

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How do DBs work?

```
sqlite>insert into student
  values (123, "Smith");
sqlite>select * from
  student;
```

| student | |
|---------|-------|
| ssn | name |
| 123 | Smith |

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```
create table student (ssn fixed, name char(20));
insert into student values(123, "Smith");
insert into student values(234, "Tompson");
insert into student values(345, "Atkinson");

-- see what we have inserted
select * from student;
```

| ssn | name |
|-------|----------|
| ----- | |
| 123 | Smith |
| 234 | Tompson |
| 345 | Atkinson |

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How do DBs work?

```
sqlite>create table takes (
  ssn fixed,
  cid char(10),
  grade fixed);
```

| takes | | |
|-------|-----|-------|
| ssn | cid | grade |
| | | |

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```
-- register students in classes and give them grades

drop table if exists takes;
create table takes (ssn fixed, cid char(10), grade fixed);

insert into takes values( 123, "db", 4);
insert into takes values( 123, "os", 3);
insert into takes values( 345, "os", 3);
insert into takes values( 345, "graphics", 4);
```

| |
|-----------------------------------|
| Smith,John,123,db,A,os,B |
| Tompson,Peter,234 |
| Atkinson,Mary,345,os,B,graphics,A |

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```
-- see what we inserted

select * from takes;
```

| ssn | cid | grade |
|-----|----------|-------|
| 123 | db | 4 |
| 123 | os | 3 |
| 345 | os | 3 |
| 345 | graphics | 4 |

| |
|-----------------------------------|
| Smith,John,123,db,A,os,B |
| Tompson,Peter,234 |
| Atkinson,Mary,345,os,B,graphics,A |

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How do DBs work - cont'd

More than one tables - joins
 Eg., roster (names only) for 'os'

| student | |
|---------|------|
| ssn | name |
| | |

| takes | | |
|-------|-----|-------|
| ssn | cid | grade |
| | | |

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How do DBs work - cont'd

```
sqlite> select name
from student, takes
where student.ssn = takes.ssn
and takes.cid = 'os'
```

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-- find the os class roster

select name from student, takes
where student.ssn = takes.ssn
and cid="os";

name

Smith
Atkinson

| |
|-----------------------------------|
| Smith,John,123,db,A,os,B |
| Tompson,Peter,234 |
| Atkinson,Mary,345,os,B,graphics,A |

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Views - a powerful tool!

what and why?

- suppose secy is allowed to see **only** ssn's and GPAs, but not individual grades
- -> VIEWS!

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Views

sqlite> create view fellowship as (
select ssn, avg(grade)
from takes group by ssn);

| takes | | |
|-------|----------|-------|
| ssn | cid | grade |
| 123 | db | 4 |
| 123 | os | 3 |
| 345 | os | 3 |
| 345 | graphics | 4 |

| ssn | avg(grade) |
|-----|------------|
| 123 | 3.5 |
| 345 | 3.5 |

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Views

Views = 'virtual tables'

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Views

sqlite> select * from fellowship;

| ssn | cid | grade |
|-----|----------|-------|
| 123 | db | 4 |
| 123 | os | 3 |
| 345 | os | 3 |
| 345 | graphics | 4 |

| ssn | avg(grade) |
|-----|------------|
| 123 | 3.5 |
| 345 | 3.5 |

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Views

sql> grant select on fellowship to secy;
(*'grant' not supported in sqlite*)

| ssn | cid | grade |
|-----|----------|-------|
| 123 | db | 4 |
| 123 | os | 3 |
| 345 | os | 3 |
| 345 | graphics | 4 |

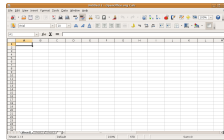

| ssn | avg(grade) |
|-----|------------|
| 123 | 3.5 |
| 345 | 3.5 |

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Iterating: advantages over (flat) files

- **logical** and **physical** data independence, because data layout, security etc info: stored **explicitly** on the disk
- concurrent access and transaction support

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Disadvantages over (flat) files?

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Disadvantages over (flat) files

- Price
- additional expertise (SQL/DBA)

hence: over-kill for small, single-user data sets

But: mobile phones (eg., android) use sqlite;
some versions of firefox do, too: ./mozilla/.../cookies.sqlite etc

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Fundamental concepts

- 3-level architecture
- logical data independence
- physical data independence

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3-level architecture

- view level
- logical level
- physical level

```

graph TD
    v1[v1] --- L[ ]
    v2[v2] --- L
    L --- P[ ]
  
```

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3-level architecture

- view level
- logical level: eg., tables
 - STUDENT(ssn, name)
 - TAKES (ssn, cid, grade)
- physical level:
 - how are these tables stored, how many bytes / attribute etc

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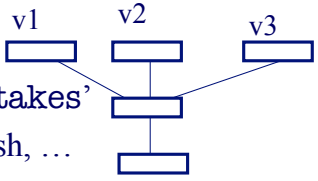
3-level architecture

- view level, eg:
 - v1: select ssn from student
 - v2: select ssn, c-id from takes
- logical level
- physical level

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3-level architecture

- view level -> 'fellowship' 
- logical level -> 'student' 'takes'
- physical level -> indices, hash, ...

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3-level architecture

- -> hence, **physical** and **logical** data independence:
- logical D.I.:
 - ???
- physical D.I.:
 - ???

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3-level architecture

- -> hence, **physical** and **logical** data independence:
- logical D.I.:
 - can add (drop) column; add/drop table
- physical D.I.:
 - can add index; change record order

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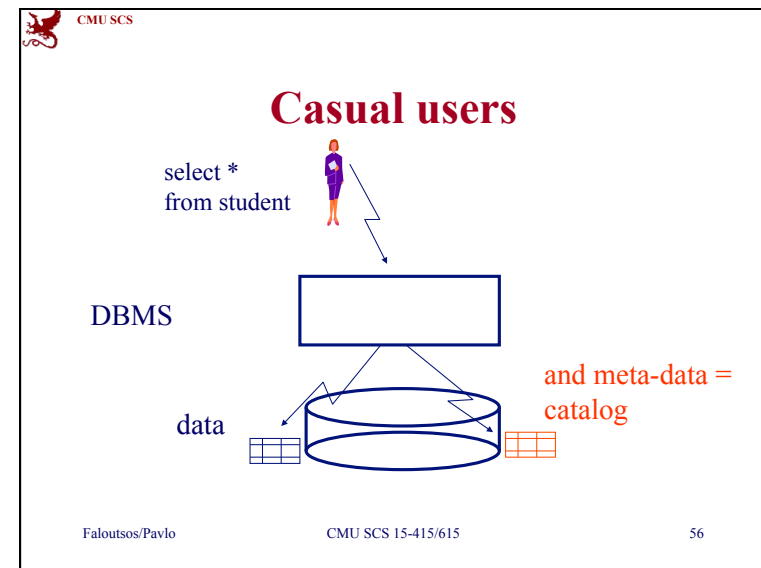
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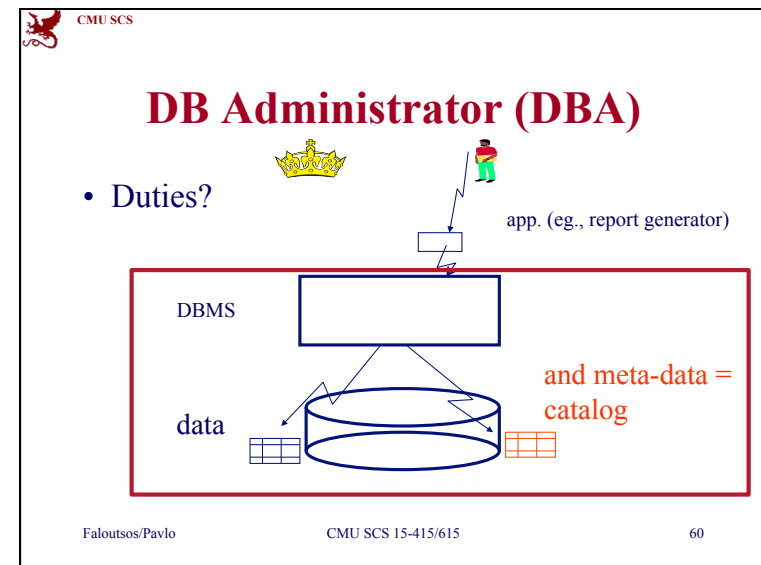
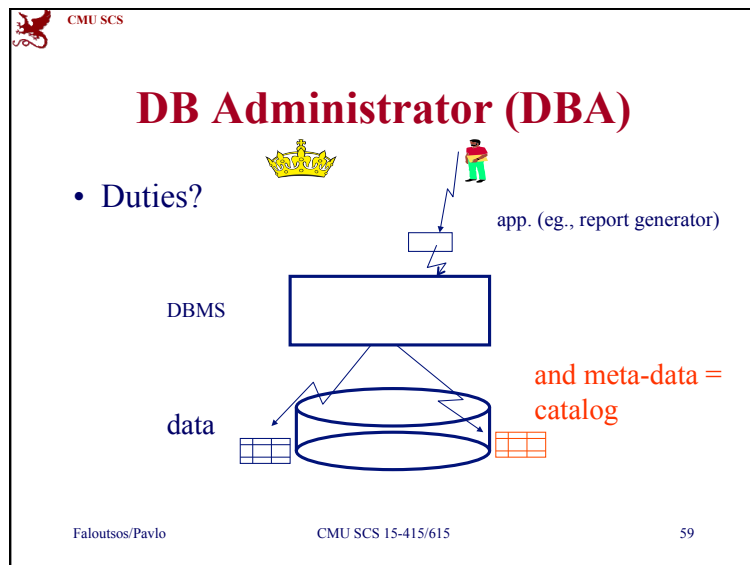
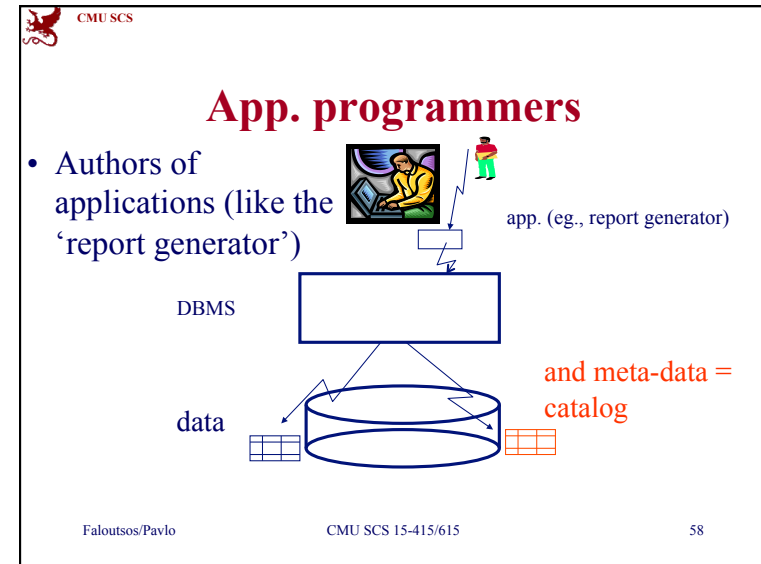
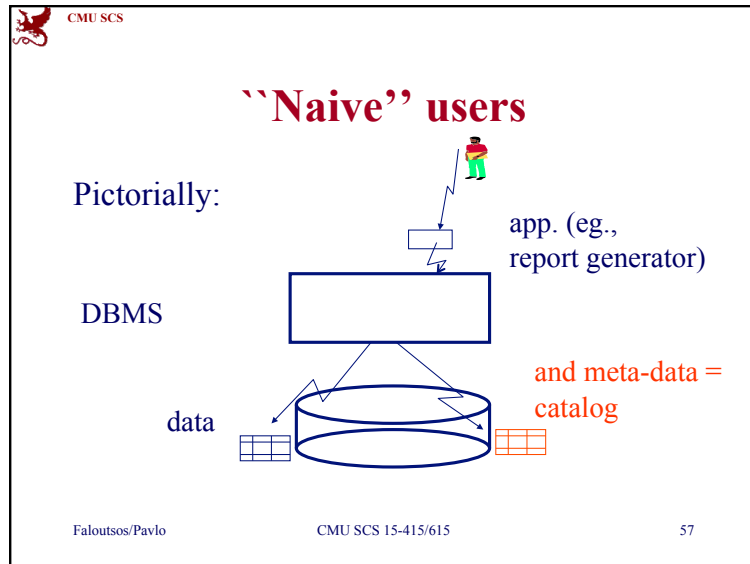
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Database users


- ‘naive’ users
- casual users
- application programmers
- [DBA (Data base administrator)]

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DB Administrator (DBA)

- schema definition ('logical' level)
- physical schema (storage structure, access methods)
- schema modifications
- granting authorizations
- integrity constraint specification

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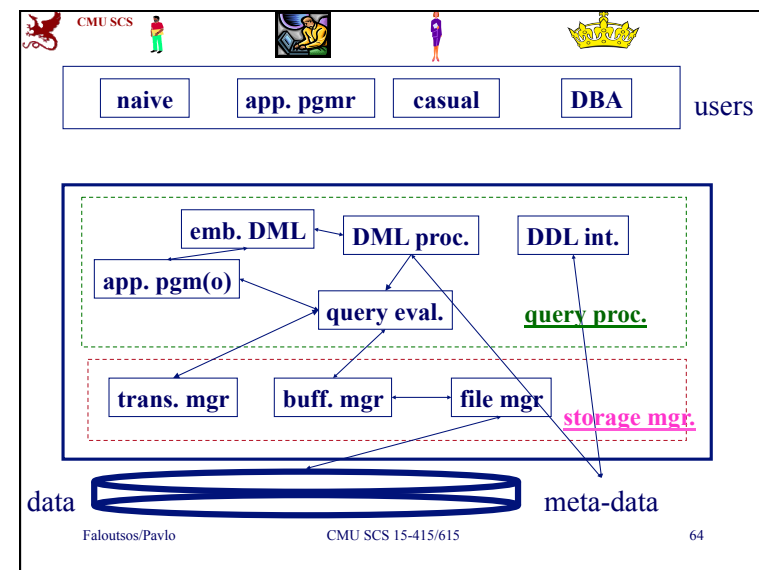
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Overall system architecture

- [Users]
- DBMS
 - query processor
 - storage manager
- [Files]

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Overall system architecture

- query processor
 - DML compiler
 - embedded DML pre-compiler
 - DDL interpreter
 - Query evaluation engine

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Overall system architecture (cont'd)

- storage manager
 - authorization and integrity manager
 - transaction manager
 - buffer manager
 - file manager

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Overall system architecture (cont'd)

- Files
 - data files
 - data dictionary = catalog (= meta-data)
 - indices
 - statistical data

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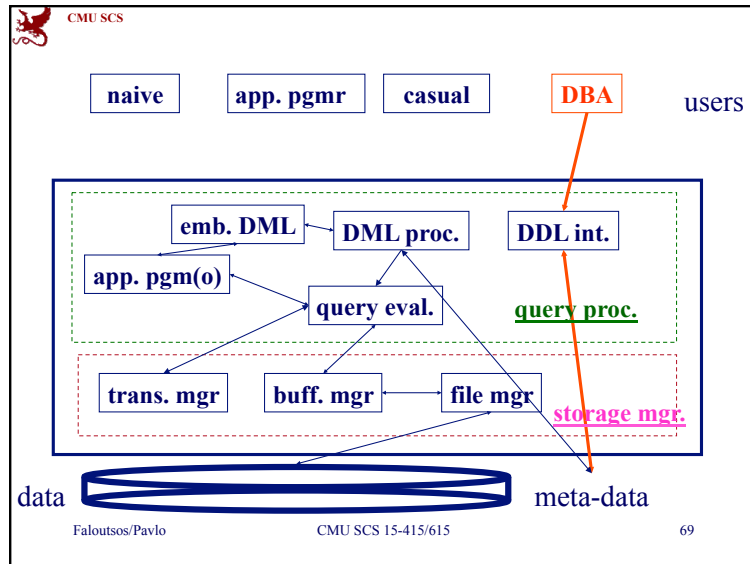


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Some examples:

- DBA doing a DDL (data definition language) operation, eg.,
create table student ...

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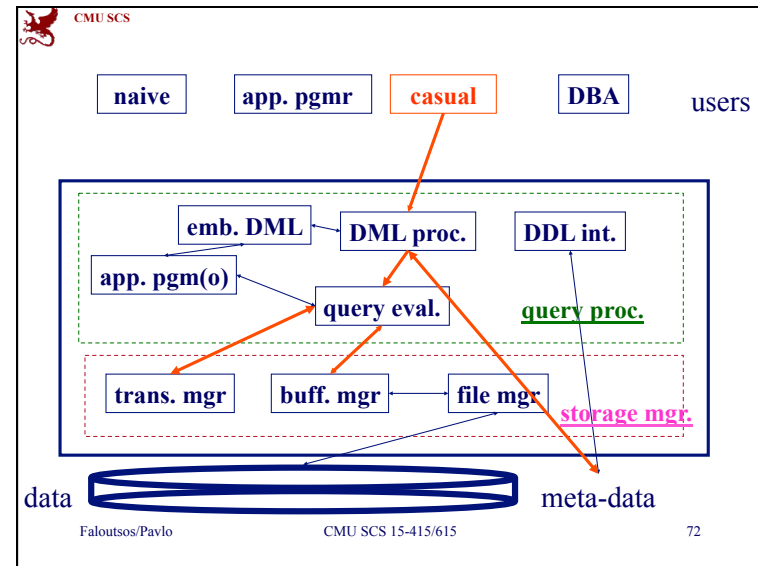
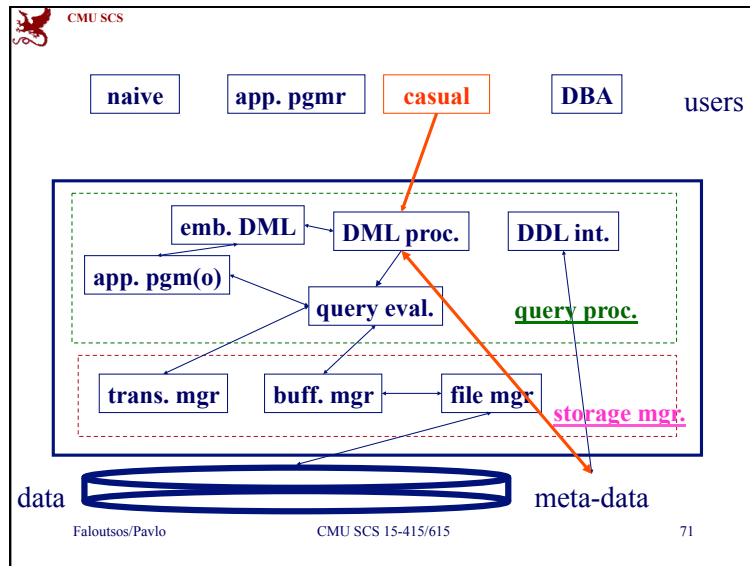


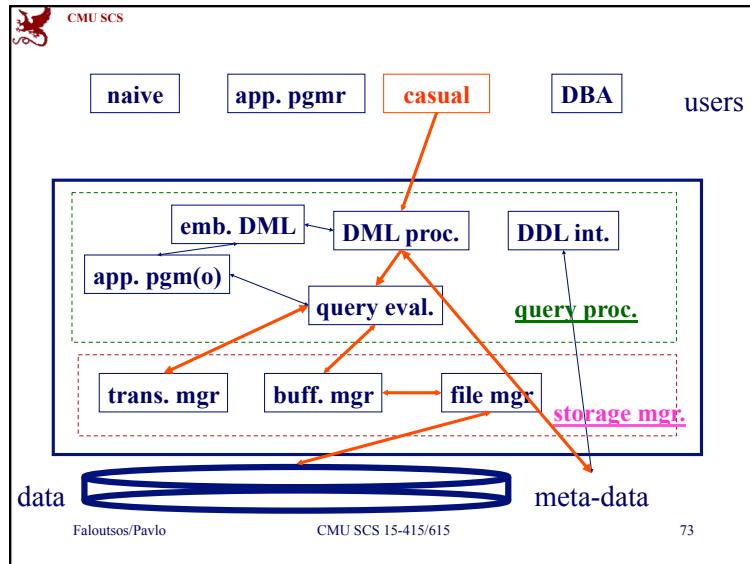
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Some examples:

- casual user, asking for an update, eg.:
 update student
 set name to 'smith'
 where ssn = '345'

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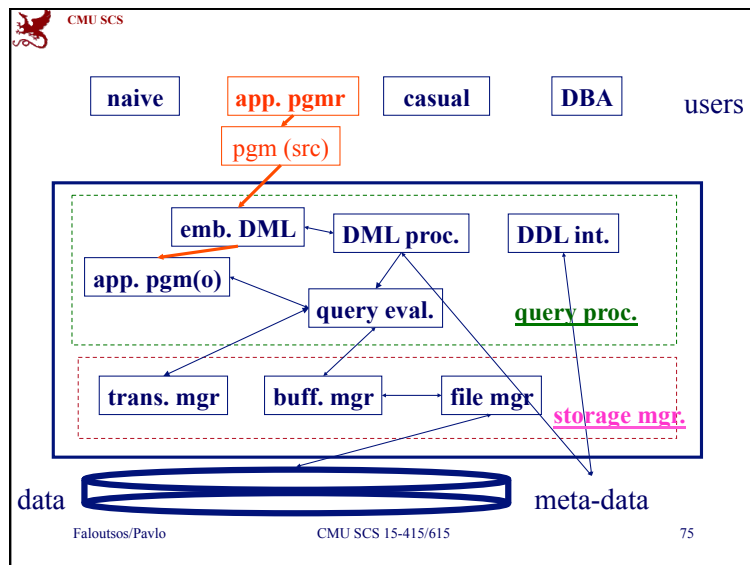


Some examples:

- app. programmer, creating a report, eg


```
main(){
  ...
  exec sql "select * from student"
  ...
}
```

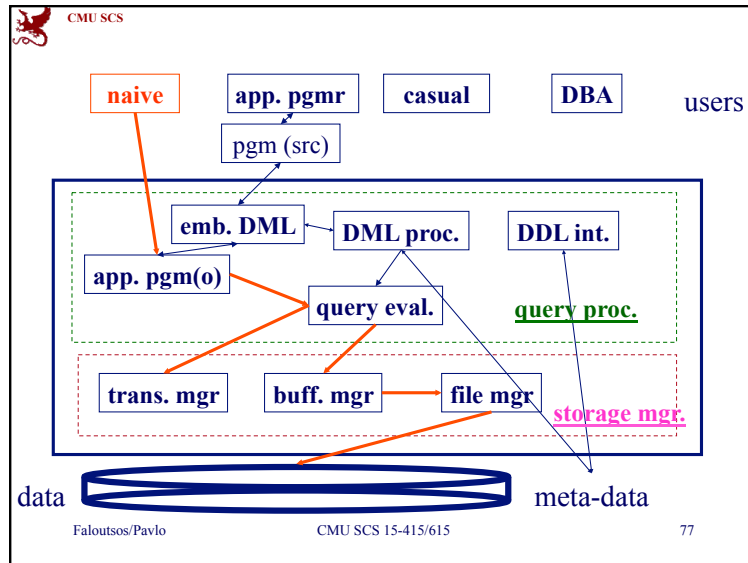
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Some examples:

- 'naive' user, running the previous app.

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Conclusions

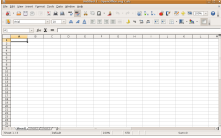

- (relational) DBMSs: electronic record keepers
- customize them with **create table** commands
- ask SQL queries to retrieve info

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Conclusions cont'd

main advantages over (flat) files & scripts:

- **logical + physical data independence** (ie., flexibility of adding new attributes, new tables and indices)
- **concurrency control and recovery**

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